

1. (a) [5 marks] (B/S) We have

$$\int_0^y x^\alpha dx = \frac{y^{\alpha+1}}{\alpha+1},$$

so $Z_p = 2^{\alpha+1}/(\alpha+1)$ and the cdf $F(x) = \frac{x^{\alpha+1}}{2^{\alpha+1}}$ for $0 \leq x \leq 2$. The inverse is $F^{-1}(u) = 2u^{1/(\alpha+1)}$ so

$$X = F^{-1}(U) = 2U^{1/(\alpha+1)}.$$

- (b) [15 marks] (i) [5 marks] (S) We have

$$M = \sup_{0 \leq x \leq 2} \frac{p(x)}{q(x)} = \frac{2^\alpha/Z_\alpha}{1/2} = \alpha + 1,$$

so the acceptance probability function

$$h(y) = \frac{p(y)}{Mq(y)} = \frac{y^\alpha}{2^\alpha}.$$

The algorithm is: (1) Simulate $Y \sim U(0, 2)$ and $U \sim U(0, 1)$; (2) if $U \leq h(Y)$ return $X = Y$ and otherwise goto (1).

- (ii) [6 marks] (S)

```
rpalpha <- function(alpha, n=1) {
  #simulates n samples X~x^alpha
  X <- numeric(n)
  for (i in 1:n) {
    FINISHED=FALSE
    while (!FINISHED) {
      y <- 2*runif(1)
      u <- runif(1)
      FINISHED <- (u <= (y^alpha/2^alpha))
    }
    X[i] <- y
  }
  return(X)
}
```

- (iii) [4 marks] (S but a little harder) If $U \sim U(0, 1)$ and $Y \sim U(0, 2)$, this is

$$M = \mathbb{P}(U \leq h(Y)) = \frac{1}{2} \int_0^2 h(y) dy = \frac{1}{2^{\alpha+1}} \int_0^2 x^\alpha dy = \frac{1}{\alpha+1}.$$

Since the number of trials, N say, has a geometric distribution with success probability ξ , we have $\mathbb{E}(N) = 1/\xi = \alpha + 1$. [Agrees with the general result $1/M$ which they know.]

- (c) [5 marks] (N) Let $x \in [0, 1]$, and calculate

$$\begin{aligned} \mathbb{P}(Y \leq x, U \leq 2Y) &= \int_0^x \left[\int_0^1 1(u \leq 2y) du \right] dy \\ &= \int_0^x \min\{1, 2y\} dy. \end{aligned}$$

Also,

$$\mathbb{P}(U \leq 2Y) = \int_0^{1/2} 2y dy + \int_{1/2}^1 1 dy = \frac{3}{4}.$$

We obtain

$$\begin{aligned}
 \mathbb{P}(X \leq x) &= \mathbb{P}(Y \leq x \mid U \leq 2Y) \\
 &= \frac{\mathbb{P}(Y \leq x, U \leq 2Y)}{\mathbb{P}(U \leq 2Y)} \\
 &= \frac{4}{3} \begin{cases} x^2, & 0 \leq x \leq 1/2 \\ 1/4 + (x - 1/2), & 1/2 \leq x \leq 1 \end{cases} \\
 &= \begin{cases} \frac{4}{3}x^2, & 0 \leq x \leq 1/2 \\ \frac{1}{3}(4x - 1), & 1/2 \leq x \leq 1. \end{cases}
 \end{aligned}$$

This is the conditional probability of X being accepted in each try, and the other tries are independent. [We went through the rejection sampling proof even more formally in the lectures, and more detailed explanation is welcome but not mandatory.]

2. (a) [6 marks] (B) Draw $Y_1, \dots, Y_n \sim q$ and

$$I_n^{(q)}(f) = \frac{1}{n} \sum_{k=1}^n f(Y_k) w(Y_k), \quad w(y) = \frac{p(y)}{q(y)} \text{ if } q(y) > 0 \text{ and } 0 \text{ otherwise.}$$

The unbiasedness comes from

$$\mathbb{E}[I_n^{(q)}(f)] = \mathbb{E}[f(Y_1)w(Y_1)] = \sum_{x \in \mathbb{X}: q(y) > 0} f(y) \frac{p(y)}{q(y)} q(y) = \sum_{x \in \mathbb{X}: q(y) > 0} f(y) p(y) = \mathbb{E}_p[f(X)],$$

because if $q(y) = 0$ then $p(y) = 0$. [This must be explained well in a perfect answer.]

- (b) [12 marks] (S)

- (i) [6 marks] Set $X_0 = 1$. The state X_k for $k = 2, 3, \dots$ is determined as follows: (1) simulate $Y_k \sim U\{1, 2, \dots, m\}$ and $U_k \sim U(0, 1)$; (2) if $U_k < \alpha(Y_k | X_k)$ set $X_k = Y_k$ and otherwise set $X_k = X_{k-1}$. In this algorithm

$$\begin{aligned} \alpha(y | x) &= \min \left\{ 1, \frac{\hat{p}(y)\hat{q}(x)}{\hat{p}(x)\hat{q}(y)} \right\} \\ &= \min \{ 1, \exp(\sqrt{x} - \sqrt{y}) \}, \end{aligned}$$

since $\hat{q}(x) = \hat{q}(y) = 1/m$ and the normalising constants $\exp(c)$ cancel.

- (ii) [6 marks] For example

```
mh <- function(n, m) {
#mcmc simulating n steps of MC targeting exp(-sqrt(x)), x=1,2,...,m
  X <- numeric(n)
  x <- 1
  for (k in 1:n) {
    y <- ceiling(runif(1)*m)
    u <- runif(1)
    r
    if (u < exp(-sqrt(y)+sqrt(x))) {
      x <- y
    }
    X[k] <- x
  }
  return(X)
}
```

- (c) [7 marks] (N)

- (i) The transition probability K can be written down as

$$\begin{aligned} [K]_{(x,z),(y,t)} &= \mathbb{P}((X_k, Z_k) = (y, t) | (X_{k-1}, Z_{k-1}) = (x, z)) \\ &= q(y | x) h(t | y) \alpha((y, t) | (x, z)) \\ &\quad + 1((y, t) = (x, z)) \rho(x, z), \end{aligned}$$

with

$$\begin{aligned} \alpha((y, t) | (x, z)) &= \min \left\{ 1, \frac{t}{z} \frac{q(x | y)}{q(y | x)} \right\}, \quad (\text{if } zq(y | x) > 0 \text{ and } 0 \text{ otherwise}), \\ \rho(x, z) &= 1 - \sum_{y \in \mathbb{X}, t \in \mathbb{N}} q(y | x) h(t | y) \alpha((y, t) | (x, z)). \end{aligned}$$

Reversibility requires us to show that $\pi(x, z)[K]_{(x, z), (y, t)} = \pi(y, t)[K]_{(y, t), (x, z)}$, which is obvious in case $(y, t) = (x, z)$, and for $(y, t) \neq (x, z)$ with $zq(y | x) > 0$ and $tq(x | y) > 0$,

$$\begin{aligned}\pi(x, z)[K]_{(x, z), (y, t)} &= h(z | x)zq(y | x)h(t | y)\alpha((y, t) | (x, z)) \\ &= h(z | x)h(t | y) \min \{zq(y | x), tq(x | y)\} \\ &= h(t | y)tq(x | y)h(z | x) \min \left\{ \frac{zq(y | x)}{tq(x | y)}, 1 \right\} \\ &= \pi(y, t)[K]_{(y, t), (x, z)}.\end{aligned}$$

[The idea of the proof is essentially the same as shown for the MH, but expressions are not exactly the same.]

- (ii) If the Markov chain is irreducible and has the invariant probability $\pi(x, z) = h(z | x)z$, then

$$\begin{aligned}\lim_{n \rightarrow \infty} \frac{1}{n} \sum_{k=1}^n f(X_k) &= \mathbb{E}_\pi[f(X, Z)], \quad \text{with } f(x, z) = f(x) \\ &= \sum_{x \in \mathbf{X}} \sum_{z \in \mathbf{N}} \pi(x, z)f(x) \\ &= \sum_{x \in \mathbf{X}} f(x) \sum_{z \in \mathbf{N}} h(z | x)z \\ &= \sum_{x \in \mathbf{X}} f(x)p(x) = \mathbb{E}_p[f(X)].\end{aligned}$$

3. Let $X = [X_1, \dots, X_p]$ be an $n \times p$ real matrix of rank p with $p < n$ column vectors X_i , $i = 1, \dots, p$, and let $\beta = (\beta_1, \beta_2, \dots, \beta_p)^T$ and $y = (y_1, y_2, \dots, y_n)^T$ be real vectors. The normal equations are $X^T X \beta = X^T y$. Consider solving the normal equations for β .

- (a) [7 marks] (i) [4 marks] (B/S close to collection question) $X^T(X\beta)$ is $2np$ operations. $(X^T X)\beta$ is $np^2 + p^2$. R uses the latter to evaluate `t(X)%*%X%*%beta`.
 (ii) [3 marks] (B) $X^T X \beta = X^T y \Leftrightarrow R^T Q^T Q R \beta = R^T Q^T y \Leftrightarrow R^T R \beta = R^T Q^T y$ (since Q is orthogonal) and hence if β satisfies $R\beta = Q^T y$ then it is a solution of the normal equations.
- (b) [6 marks] Let $X_{2:p} = [X_2, X_3, \dots, X_p]$. We have

$$\begin{aligned} X &= [X_1, X_{2:p}] \\ Q &= [Q_1, Q'] \\ X &= QR \\ [X_1, X_{2:p}] &= [Q_1, Q'] \begin{pmatrix} a & r \\ 0_{(p-1) \times 1} & R' \end{pmatrix} \\ [X_1, X_{2:p}] &= [aQ_1, Q_1 r + Q' R'] \end{aligned}$$

We can read off $X_1 = aQ_1$ and $X_{2:p} = Q_1 r + Q' R'$ from the last line. Now $|Q_1| = 1$ $a = |X_1|$ and so $Q_1 = X_1/a$. Multiplying $X_{2:p} = Q_1 r + Q' R'$ by Q_1^T gives $r = Q_1^T X_{2:p}$. Finally, if $X' = X_{2:p} - Q_1 r$ then $X' = Q' R'$.

Now R' is upper triangular with positive diagonal entries (since R was) and Q' is clearly orthogonal, hence $Q' R'$ is the QR-factorisation of X' .

- (c) [4 marks] Each time we apply this we reduce the number of columns of X by one. Our Algorithm is
- Step 1. If $p = 1$ then $Q = X/|X|$ and $R = |X|$ and we are done.
 Step 2. Otherwise (if $p > 1$) then
 Step 2.1 set $Q_1 = X_1/|X_1|$, $r = Q_1^T X_{2:p}$ and $X' = X_{2:p} - Q_1 r$.
 Step 2.2 call this algorithm to compute the QR factorisation $X' = Q' R'$.
 Step 2.3 Assemble Q and R from Q_1, Q', r and R' .

- (d) [8 marks] S/N

```
my.qr<-function(X) {
  norm<-function(v) {sqrt(sum(v^2))}

  n=dim(X)[1]; p=dim(X)[2] #assume p<n

  R=matrix(NA,p,p)
  Q=matrix(NA,n,p)

  R[1,1]=norm(X[,1])
  Q[,1]=X[,1]/R[1,1]

  if (p==1) return(list(Q=Q,R=R))

  R[1,2:p]=t(Q[,1])%*%X[,2:p]
  R[2:p,1]=0
  Xp=X[,2:p,drop=F]-Q[,1,drop=F]%*%R[1,2:p,drop=F]
```

```
QR=my.qr(Xp)

Q[,2:p]=QR$Q
R[2:p,2:p]=QR$R

return(list(Q=Q,R=R))
}
```