

# MiRitH

(MinRank in the Head)

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**PQC Workshop Oxford** 

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Cryptography Research Center





EUF-CMA secure in the ROM assuming hardness of MinRank



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2. Approach MPC-in-the-Head:



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#### 2. Approach <u>MPC-in-the-Head</u>:

- 1. MPC protocol to verify a shared solution of MinRank
- 2. Zero-Knowledge proof of a solution
- 3. Signature scheme from Fiat-Shamir transform



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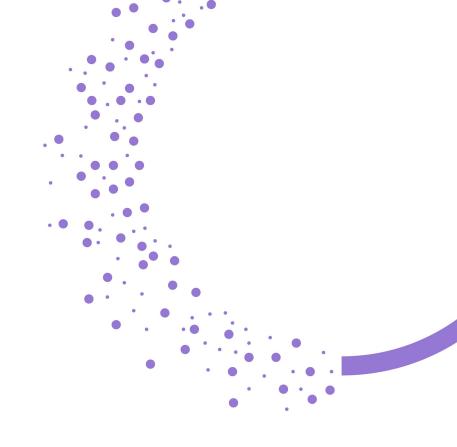
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#### 3. Parameters, Security and Performance



# The MinRank problem



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Given: An integer 
$$r$$
, and  $k + 1$  matrices  $M_0, M_1, \dots, M_k \in \mathbb{F}_q^{m \times n}$ 

**<u>Find:</u>**  $\alpha_1, \dots, \alpha_k \in \mathbb{F}_q$  such that  $E = M_0 + \sum_{i=1}^k \alpha_i M_i$  has  $rank(E) \le r$ 

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MinRank as decoding problem:

Given: An integer r, and 
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MinRank as decoding problem:

Gen. Matrix 
$$G = \begin{pmatrix} Vec(M_1) \\ \vdots \\ Vec(M_k) \end{pmatrix} \in \mathbb{F}_q^{k \times (n \cdot m)}$$

 $Vec(M_0) = (\alpha_1, ..., \alpha_k) \cdot G + Vec(E)$ , where  $rank(E) \le r$ 

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#### Type of instances we use:

Given: An integer r, and 
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Random matrices

Type of instances we use:

- Random secret
- Random *E*



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Starting point *N*-Party MPC protocol

Starting point *N*-Party MPC protocol Given: function f, value z, share  $x_i$  of xGoal: Verify if f(x) = z, with  $x = \sum x_i$ Output: accept :  $P'_i$ 's think they **do** share x. reject :  $P'_i$ 's think they **do not** share x

Starting point *N*-Party MPC protocol

 $\begin{array}{cccc} P_1 & \longleftrightarrow & P_2 \\ x_1 & & & x_2 \end{array} \\ \uparrow & \boxtimes & \uparrow \\ P_3 & \longleftrightarrow & P_N \\ x_3 & & & x_N \end{array}$ 

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Starting point *N*-Party MPC protocol

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Given: function f, value z, share  $x_i$  of xGoal: Verify if f(x) = z, with  $x = \sum x_i$ Output:  $accept : P'_i$ s think they **do** share x.  $reject : P'_i$ s think they **do** not share x

False-Positive-Rate =  $\Pr[\text{accept} | f(x) \neq z]$ 

No information on  $x_i$  leaked to  $P_j$  for  $j \neq i$ 



Given: MPC protocol

## <u>Goal:</u> zero-knowledge proof of knowledge

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(Prover *P* wants to proof knowledge of *x* with f(x) = z to *V*)

#### **Prover**

Given: MPC protocol

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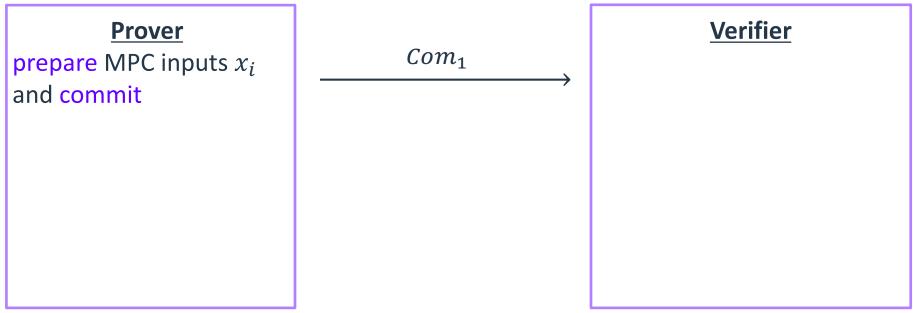
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#### **Prover**



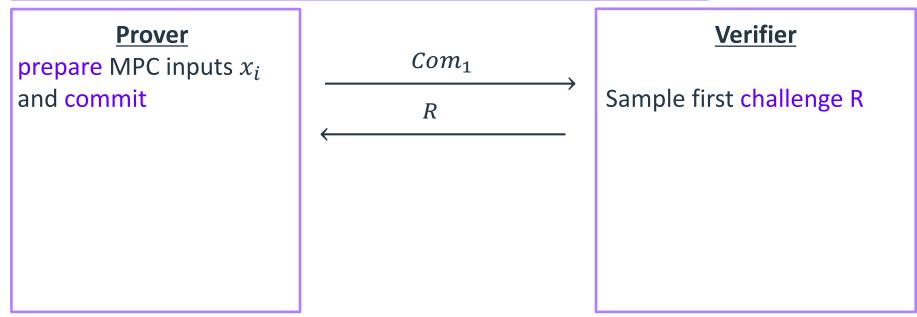


## <u>Goal:</u> zero-knowledge proof of knowledge



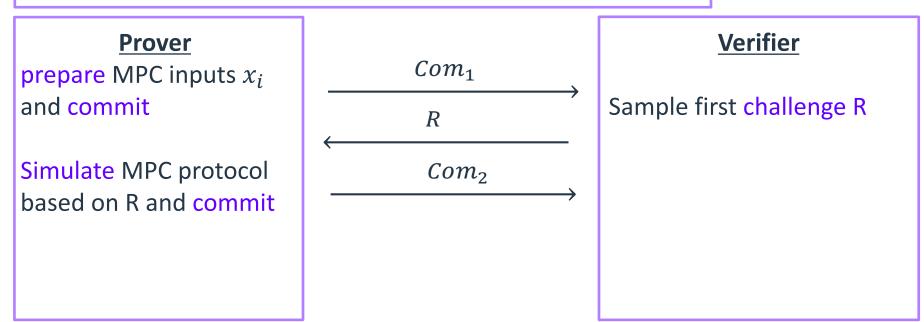


## <u>Goal:</u> zero-knowledge proof of knowledge



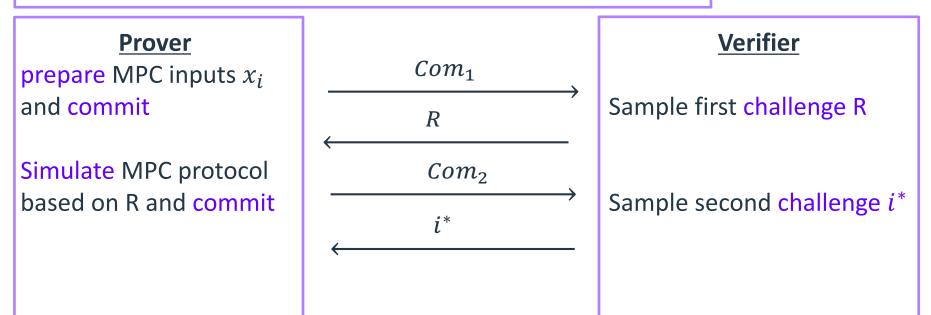
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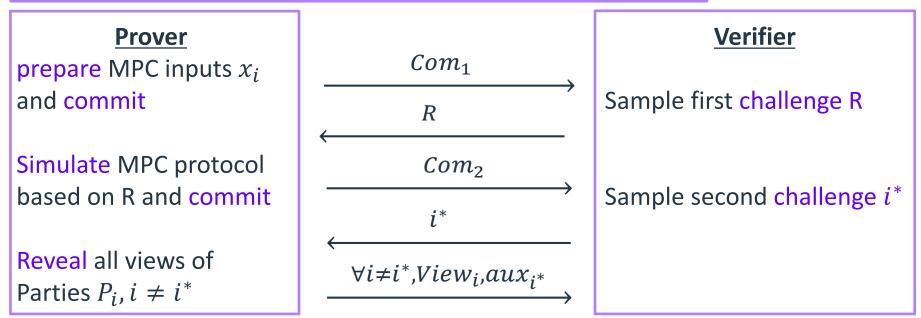
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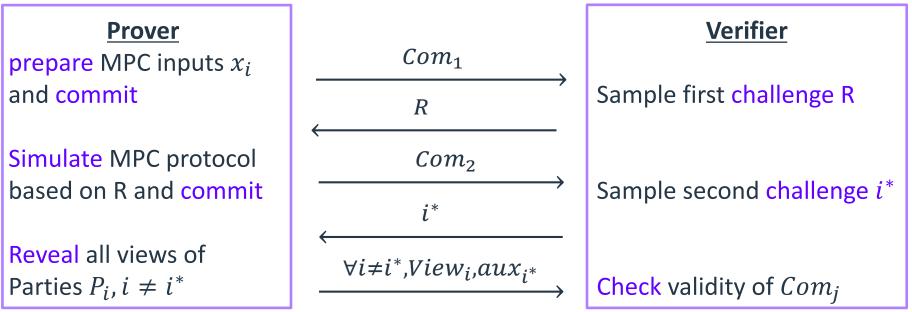
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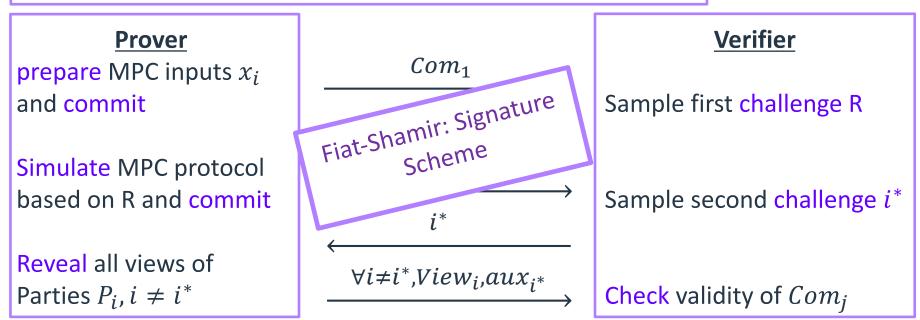
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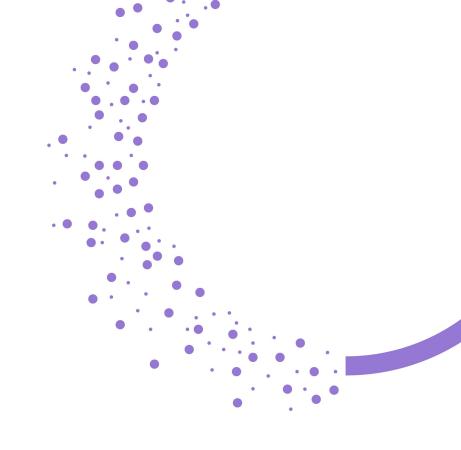
Given: MPC protocol

# <u>Goal:</u> zero-knowledge proof of knowledge





# **Design rationale**



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Models MinRank as a bilinear system

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$$\left(M_0 + \sum_{i=1}^k \frac{\beta_i M_i}{K}\right) \cdot \binom{l_{n-r}}{K} = 0$$

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Solving system  $\Rightarrow$  Solving MinRank!

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$$\beta_i = \alpha_i$$

Models MinRank as a bilinear system

$$\left(M_0 + \sum_{i=1}^k \alpha_i M_i\right) \cdot \binom{I_{n-r}}{K} = 0$$

Solving system  $\Rightarrow$  Solving MinRank!

## **Kipnis-Shamir modelling**

Models MinRank as a bilinear system

$$\left( M_{0} + \sum_{i=1}^{k} \alpha_{i} M_{i} \right) \cdot {\binom{l_{n-r}}{K}} = 0$$
Solving
$$M_{\vec{\alpha}} \cdot {\binom{l_{n-r}}{K}} = 0 \iff M_{\vec{\alpha}}^{L} = -M_{\vec{\alpha}}^{R} \cdot K$$

Solving system  $\Rightarrow$  Solving MinRank!

### Kipnis-Shamir modelling

Models MinRank as a bilinear system

$$\begin{pmatrix} M_0 + \sum_{i=1}^{k} \alpha_i M_i \end{pmatrix} \cdot \begin{pmatrix} I_{n-r} \\ K \end{pmatrix} = 0 \qquad \text{Solving syst}$$

$$M_{\vec{\alpha}} \cdot \begin{pmatrix} I_{n-r} \\ K \end{pmatrix} = 0 \iff M_{\vec{\alpha}}^L = -M_{\vec{\alpha}}^R \cdot K$$

Solving system  $\Rightarrow$  Solving MinRank!

Knowledge of MinRank solution  $\vec{\alpha}$   $\Leftrightarrow$ Knowledge of K such that  $M_{\vec{\alpha}}^L = -M_{\vec{\alpha}}^R \cdot K$ 

 $\vec{\alpha}$  solution of MinRank problem  $M_0, M_1, \dots, M_k$ 

 $\vec{\alpha} = \sum_{i=1}^{N} \vec{\alpha}_i$  and  $K = \sum_{i=1}^{N} K_i$ 

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$$\begin{array}{cccc}
P_1 & \longleftrightarrow & P_2 \\
(\vec{\alpha}_1, K_1) & & (\vec{\alpha}_2, K_2)
\end{array}$$

$$\begin{array}{cccc}
\uparrow & \swarrow & \uparrow \\
P_3 & & P_N \\
(\vec{\alpha}_3, K_3) & \longleftrightarrow & (\vec{\alpha}_N, K_N)
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\end{array}$ 

 $\vec{\alpha} = \sum_{i=1}^{N} \vec{\alpha}_i$  and  $K = \sum_{i=1}^{N} K_i$ 

Goal: Verify parties share  $(\vec{\alpha}, K)$  s.t.  $M_{\vec{\alpha}}^L = -M_{\vec{\alpha}}^R \cdot K$ Output: <u>accept</u> :  $P'_i$ s think they **do** share  $(\vec{\alpha}, K)$ <u>reject</u> :  $P'_i$ s think they **don't** share  $(\vec{\alpha}, K)$ 

No information on  $(\vec{\alpha}_i, K_i)$  leaked

 $\vec{\alpha}$  solution of MinRank problem  $M_0, M_1, \dots, M_k$ 

 $\vec{\alpha} = \sum_{i=1}^{N} \vec{\alpha}_i$  and  $K = \sum_{i=1}^{N} K_i$  $(\vec{\alpha}_1, K_1)$ MiRitH: MPC verifies (X, Y, Z) satisfies  $X \cdot Y = Z$ 

 $\longleftrightarrow \begin{array}{c} P_N \\ (\vec{\alpha}_N, K_N) \end{array}$ 

<u>Goal</u>: Verify parties share  $(\vec{\alpha}, K)$  s.t.  $M_{\overrightarrow{\alpha}}^L = -M_{\overrightarrow{\alpha}}^R \cdot K$ 

#### Output:

<u>accept</u> :  $P'_i$ 's think they **do** share  $(\vec{\alpha}, K)$ 

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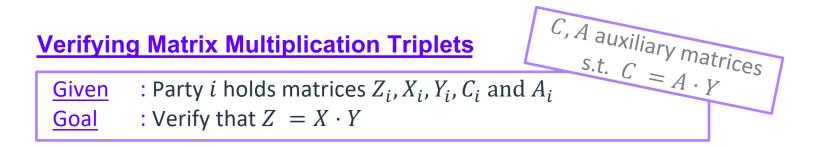
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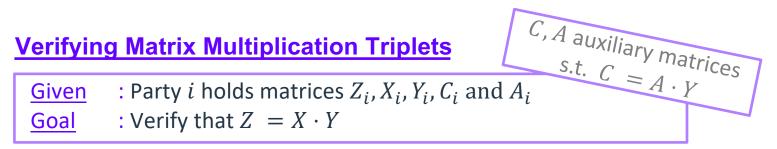
### **Verifying Matrix Multiplication Triplets**

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Given : Party *i* holds matrices 
$$Z_i, X_i, Y_i, C_i$$
 and  $A_i$ 

**<u>Goal</u>** : Verify that  $Z = X \cdot Y$ 





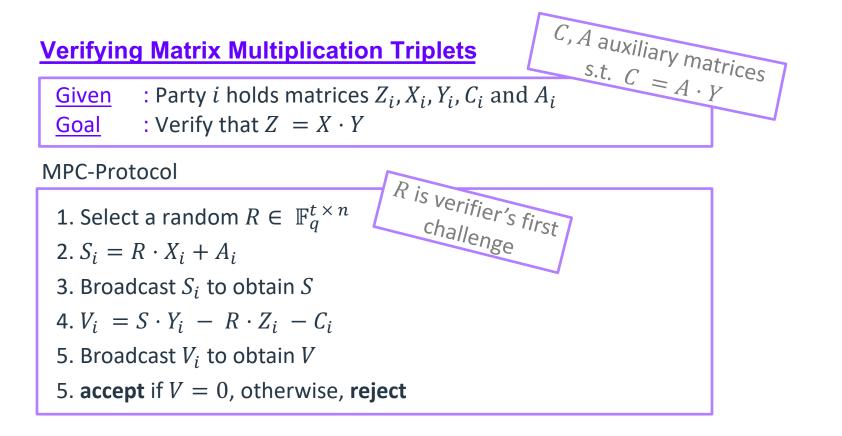
MPC-Protocol

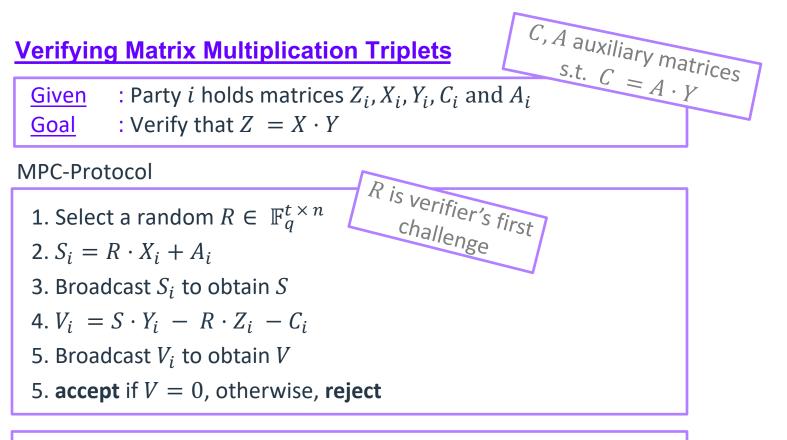
# **Verifying Matrix Multiplication Triplets**

- C, A auxiliary matrices s.t.  $C = A \cdot Y$ : Party *i* holds matrices  $Z_i, X_i, Y_i, C_i$  and  $A_i$ Given
- : Verify that  $Z = X \cdot Y$ Goal

#### MPC-Protocol

- 1. Select a random  $R \in \mathbb{F}_q^{t \times n}$
- 2.  $S_i = R \cdot X_i + A_i$
- 3. Broadcast  $S_i$  to obtain S
- 4.  $V_i = S \cdot Y_i R \cdot Z_i C_i$
- 5. Broadcast  $V_i$  to obtain V
- 5. accept if V = 0, otherwise, reject





<u>Correctness</u> : If  $Z = X \cdot Y$  and  $C = A \cdot Y$ , then parties **accept** <u>False-Positive rate</u>: If not, the Parties **accept** with prob.  $q^{-t}$ 





1. <u>Kinpis-Shamir</u> modelling



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- 2. MPC for <u>matrix-triplet verification</u>



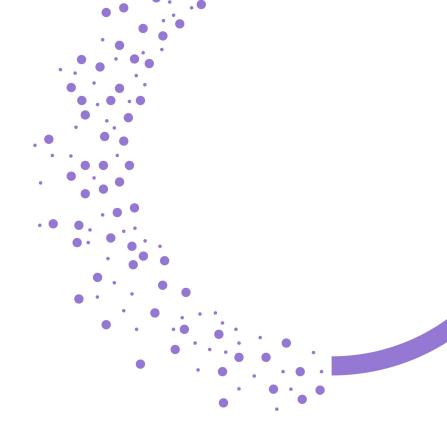
- 1. <u>Kinpis-Shamir</u> modelling
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- 2. MPC for matrix-triplet verification
- 3. <u>MPC-in-the-Head</u> (incl. hypercube, seedtrees, etc.)
- 4. Fiat-Shamir transform



## **Security, Parameters and Performance**



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Find  $\alpha \in \mathbb{F}_q^k$  such that:  $E \coloneqq M_0 + \sum_{i=1}^k \alpha_i M_i \in \mathbb{F}_q^{m \times n}$ , and  $rank(E) \leq r$ 

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1. Kernel Search (combinatorial) : Guess vectors in kernel(E)

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- 3. Big-k (combinatorial) : Guess entries of *E*

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Hybrid approach: Guess some of the  $\alpha_i$ 's, and some vectors in kernel(E)

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- 2. Support-Minors (algebraic) : Model as bilinear system of equations
- 3. Big-k (combinatorial) : Guess entries of *E*

Hybrid approach: Guess some of the  $\alpha_i$ 's, and some vectors in kernel(E) $\rightarrow$  MinRank instance of smaller dimension **Category I MinRank Parameters** 

### **Category I MinRank Parameters**

Category	set	q	m = n	k	r
I.	а	16	15	78	6
I	b	16	16	142	4

Parameters of the underlying MinRank instance

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I	а	16	15	78	6
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Parameters of the underlying MinRank instance

Category	set	Kernel-Search	Support Minors	Big-k
I	а	151	144	154
I	b	159	165	226

Complexity estimates for proposed parameters, with linear algebra constant equal to 3 in KS and Big-k, equal 2.81 for Strassen in SM.



# **Performance**

Category I a	Sig. Size	Pk size	Key Gen.	Sign	Verify
Short	5.7 kB		F2 000	~23 MCycles	
Fast	7.7 kB	129 Bytes	~53.000	~ 3	MCycles

# **Performance**

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Category I b	Sig. Size	Pk size	Key Gen.	Sign	Verify
Short	6.3 kB	129 Bytes	F2 000	~24 MCycles	
Fast	8.8 kB		~53.000	~ 4	MCycles

#### **Performance**

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Short	5.7 kB	100 D .	50.000	~23 MCycles	
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